# **Computer Game Playing**



# Deep Blue Beats World Champion (1997)



# "Game Over" — Checkers Solved in 2007



# Computer Go



# General Game Playing

General Game Players are systems

- able to accept a formal description of arbitrary games
- able to use such descriptions to play the games effectively

Cognitive Information Processing Technologies for GGP systems:

- Knowledge representation
- Reasoning
- Learning
- Rational behaviour

Unlike specialised game players (e.g. Deep Blue), they do not use algorithms designed in advance for specific games.

# Tic-Tac-Toe



# **Bidding Tic-Tac-Toe**



x	0	
3 *	3	

Chess



# **Kriegspiel**





#### Single-Player Games



# Monopoly



Poker



# **International Activities**

Websites – http://games.stanford.edu

www.general-game-playing.de

- Games
- Game Manager
- Reference Players
- Development Tools
- Literature

Annual World Cup

- First: Pittsburgh 2005
- Most recent: Barcelona 2011
- Next: Toronto 2012

German Open, Berlin 2011

# General Game Playing and Al

Why games?

- Many social, biological, political, and economic processes can be formalised as (multi-agent) games.
- General game-players are rational agents that can adapt to radically different environments without human intervention.













# Knowledge Representation for GGP:

# **Describing Games**

#### Games as State Machines



### Initial State, Terminal States, Simultaneous Moves



# Game Model (Perfect-Information Games)

An *n*-player game with perfect information is a structure with components:

 $\{r_1, ..., r_n\} - players$ 

S – set of states

 $A_1, ..., A_n - n$  sets of actions, one for each player

 $I_1, ..., I_n$  – where  $I_i \subseteq A_i \times S$ , the legality relations

u:  $S \times A_1 \times ... \times A_n \rightarrow S$  – update function

 $s_1 \in S$  – initial game state

 $t \subseteq S$  – the terminal states

 $g_1, \dots g_n$  – where  $g_i \subseteq S \times IN$ , the goal relations

### **Encoding Alternatives**

**State Machines**. Astronomically large state spaces, e.g. ~ 5000 states in Tic-Tac-Toe, ~10<sup>44</sup> states in Chess.

- Lists and Tables. Still the same size. Just switching to database states does not decrease the size of direct representation.
- **Programs**. One possibility is to write a program to generate legal moves and successor states and to evaluate goals and termination. However, which language? Java, C? What if a player wants to reason about the structure of a game in general? This is difficult if the game is encoded in procedural form.
- **Logic**. There are existing interpreters / compilers. Logic is easier to use for analysis than procedural encodings.

## **Formal Game Descriptions**

Whatever form is used, the description must give all information necessary to determine legality of moves, state transition, termination, and goals.

Nothing is assumed except for logic.

- No arithmetics
- No physics
- No common sense

(To emphasise this, game descriptions can be written in terms of nonsense symbols.)

### Game Description Language

In the Game Description Language (GDL), a game is a logic program.

GDL uses the constants 0, 1, ..., 100 and the following predicates as keywords.

- role(r) means that r is a role (i.e. a player) in the game
- init(f) means that f is true in the initial position (state)
- true(f) means that f is true in the current state
- does(r,a) means that role r does action a in the current state
- next(f) means that f is true in the next state
- legal(r,a) means that it is legal for r to play a in the current state
  - goal(r,v) means that r gets goal value v in the current state
- terminal means that the current state is a terminal state
- distinct(s,t) means that terms s and t are syntactically different

#### Tic Tac Toe



- cell(1,1,x)
- cell(1,2,b)
- cell(1,3,b)
- cell(2,1,b)
- cell(2,2,0)
- cell(2,3,b)
- cell(3,1,b)
- cell(3,2,b)
- cell(3,3,x)
- control(oplayer)

# **Bidding Tic Tac Toe**

- cell(1,1,b)
- cell(1,2,b)
- cell(1,3,b)
- cell(2,1,b)
- cell(2,2,b)
- cell(2,3,b)
- cell(3,1,b)
- cell(3, 2, b)
- cell(3,3,b)
- coins(xplayer,3)
- coins(oplayer,3)
- tiebreaker(xplayer)
- bidding\_stage







# Bidding Tic Tac Toe: Vocabulary

Object constants
 xplayer, oplayer Players
 x, o, b Marks
 bidding\_stage
 tiebreaker,
 no\_tiebreaker,
 noop

#### Functions

cell(number,number,mark), control(player), coins(player,number), tiebreaker(player) Fluents mark(number,number), bid(number,flag) Moves

Domain predicates

### **Players and Initial State**

```
role(xplayer)
```

```
role(oplayer)
```

**init**(cell(1,1,b))

init(cell(1,2,b))

init(cell(1,3,b))

init(cell(2,1,b))

init(cell(2,2,b))

init(cell(2,3,b))

init(cell(3,1,b))

init(cell(3,2,b))

init(cell(3,3,b))

init(coins(P,3)) <= role(P)</pre>

init(tiebreaker(xplayer))

init(bidding\_stage)

## Move Generator (1)

legal(P,mark(M,N)) <=
 true(cell(M,N,b)) ^
 true(control(P))</pre>

legal(xplayer,noop) <=
 true(control(oplayer))</pre>

legal(oplayer,noop) <=
 true(control(xplayer))</pre>

**Conclusions**: legal(xplayer,noop) legal(oplayer,mark(1,2)) ...

legal(oplayer,mark(3,2))



cell(1,1,x) cell(1,2,b) cell(1,3,b) cell(2,1,b) cell(2,2,o) cell(2,2,o) cell(2,3,b) cell(3,1,b) cell(3,1,b) cell(3,2,b) cell(3,2,b) cell(3,3,x) coins(xplayer,4) coins(oplayer,2) tiebreaker(oplayer) control(oplayer)

# Move Generator (2)

```
legal(P,bid(B,TB)) <=
    true(bidding_stage) ^
    true(coins(P,C)) ^
    less_or_eq(B,C) ^
    tiebreak_bid(P,TB)
tiebreak_bid(P,with_tiebreaker) <=
    true(tiebreaker(P))
tiebreak_bid(P,no_tiebreaker) <=
    role(P)</pre>
```

Conclusions:

legal(xplayer,bid(1,with\_tiebreaker))

```
legal(xplayer,bid(1,no_tiebreaker))
```

• • •

legal(oplayer,bid(1,no\_tiebreaker)



oplayer

mark(1,3)

#### Physics: Example

cell(1,1,x) cell(1,2,b) cell(1,3,b) cell(2,1,b) cell(2,2,o) cell(2,3,b) cell(3,1,b) cell(3,1,b) cell(3,2,b) cell(3,2,b) cell(3,3,x) coins(xplayer,4) coins(oplayer,2) tiebreaker(oplayer) control(oplayer) cell(1,1,x)
cell(1,2,b)
cell(1,3,o)
cell(2,1,b)
cell(2,2,o)
cell(2,3,b)
cell(3,1,b)
cell(3,2,b)
cell(3,2,b)
cell(3,3,x)
coins(xplayer,4)
coins(oplayer,2)
tiebreaker(oplayer)
bidding\_stage



### Physics (1)

next(cell(M,N,x)) <= does(xplayer,mark(M,N))</pre>

next(cell(M,N,o)) <= does(oplayer,mark(M,N))</pre>

next(bidding\_stage) <= true(control(P))</pre>

# Physics (2)

<pre>next(coins(P,C)) &lt;= true(bidding_stage) ^     winner(P) ^     does(P,bid(B,TB)) ^     true(coins(P,C1)) ^ add(C,B,C1)</pre>
<pre>next(coins(xplayer,C)) &lt;= true(bidding_stage) ^     winner(oplayer) ^     does(oplayer,bid(B,TB)) ^     true(coins(P,C1)) ^ add(C1,B,C)</pre>
<pre>next(coins(oplayer,C)) &lt;= true(bidding_stage) ^</pre>

# Physics (3)

<pre>next(tiebreaker(P))</pre>	<=	<pre>true(bidding_stage) ^ role(P) ^ winner(Q) ^ distinct(P,Q) ^ does(Q,bid(B,with_tiebreaker))</pre>
<b>next</b> (tiebreaker(P))	<=	<pre>true(bidding_stage) ^ true(tiebreaker(P)) ^ winner(Q) ^ does(Q,bid(B,no_tiebreaker))</pre>

next(cell(M,N,W)) <= true(cell(M,N,W)) ^ true(bidding\_stage)
next(control(P)) <= winner(P) ^ true(bidding\_stage)</pre>

#### **Termination and Goal Values**

```
terminal <=</pre>
```

```
line(x) \lor line(o)
```

terminal <=</pre>

¬open

line(W) <=</pre>

row(M,W) ∨

column(N,W) ∨

diagonal(W)

open <=

true(cell(M,N,b))

```
goal(xplayer,100) <= line(x)
goal(xplayer, 50) <= draw
goal(xplayer, 0) <= line(o)
goal(oplayer,100) <= line(o)
goal(oplayer, 50) <= draw
goal(oplayer, 0) <= line(x)
draw <=</pre>
```

```
\negline(x) \land \negline(o) \land \negopen
```

# Supporting Concepts

```
diagonal(W) <=
row(M,W) <=
                               true(cell(1,1,W)) \land
   true(cell(M,1,W)) ∧
   true(cell(M,2,W)) ∧
                               true(cell(2,2,₩)) ∧
                               true(cell(3,3,W))
   true(cell(M,3,W))
                            diagonal(W) <=</pre>
column(N,W) <=
                                true(cell(1,3,W)) ∧
   true(cell(1,N,W)) ∧
                                true(cell(2,2,₩)) ∧
   true(cell(2,N,W)) ∧
                                true(cell(3,1,W))
   true(cell(3,N,W))
```

### More Supporting Concepts

```
winner(P) <= does(P,bid(B1,TB1)) \land does(Q,bid(B2,TB2)) \land
              distinct(P,O) \land greater(B1,B2)
winner(P) <= does(P, bid(B, with tiebreaker)) ^
              does(O,bid(B,no tiebreaker))
winner(P) <= does(P,bid(B,no tiebreaker)) </pre>
              does(0,bid(B,no tiebreaker)) ∧
              distinct(P,O) ∧
              true(tiebreaker(0))
succ(0,1) succ(1,2) succ(2,3) ...
qreater(X,Y) <= ...</pre>
less or equal(X,Y) <= ...</pre>
add(X,Y,Z) <= \ldots
```

#### Completeness

Of necessity, game descriptions are logically incomplete in that they do not uniquely specify the moves of the players.

Every game description contains *complete definitions* for *legality, termination, goalhood,* and *update* in terms of the relations true and does.

The upshot is that in every state every player can determine legality, termination, goalhood, and—given a joint move—can update the state.

### Syntactic Restrictions

- 1. role as head of clause only appears in facts (i.e., clauses with empty body)
- 2. init only appears as head of clauses and does not depend on any of true, legal, does, next, terminal, goal
- 3. true only appears in bodies of clauses
- 4. does only appears in clause bodies, and none of legal, terminal, goal depends on does
- 5. next only appears as head of clauses

# Guaranteeing Decidability (1): Safety

A clause is <u>safe</u> if and only if every variable in the clause appears in some positive subgoal in the body.

Safe Rule:

 $r(X,Y) <= p(X,Y) \land \neg q(X,Y)$ 

Unsafe Rule:

 $r(X,Z) <= p(X,Y) \land \neg q(Y,Z)$ 

In GDL, all rules are required to be safe.

(Note that this implies all facts to be variable-free.)

### **Dependency Graph**

The dependency graph for a set of clauses is a directed graph in which

- the nodes are the relations mentioned in the head and bodies of the clauses
- there is an arc from a node p to a node q whenever p occurs in the body of a clause in which q is in the head.

r(X,Y) <= p(X,Y) ∧ q(X,Y) s(X,Y) <= r(X,Y) s(X,Z) <= r(X,Y) ∧ t(Y,Z) t(X,Z) <= s(X,Y) ∧ s(Y,X)



A set of clauses is <u>recursive</u> if its dependency graph contains a cycle. Otherwise, it is <u>non-recursive</u>.

# Guaranteeing Decidability (2): Stratification

A set of rules is said to be <u>stratified</u> if there is no recursive cycle in the dependency graph involving a negation.

Stratified Negation:

$$t(X,Y) <= q(X,Y) \land \neg r(X,Y)$$
  
r(X,Z) <= p(X,Y)  
r(X,Z) <= r(X,Y) \land r(Y,Z)

Negation that is not stratified:

r(X,Z) <= p(X,Y) r(X,Z) <= p(X,Y) ∧ ¬r(Y,Z)

In GDL, all sets of rules are required to be stratified.

# Guaranteeing Decidability (3)

If a set of rules contains a clause

$$p(s_1, \ldots, s_m) \leq b_1(\vec{t_1}) \wedge \ldots \wedge q(v_1, \ldots, v_k) \wedge \ldots \wedge b_n(\vec{t_n})$$

where p and q occur in a cycle in the dependency graph, then for every  $i \in \{1, ..., k\}$ 

- *v<sub>i</sub>* is variable-free, or
- *v<sub>i</sub>* is one of s<sub>1</sub>, ..., s<sub>m</sub>, or
- $v_i$  occurs in some  $\vec{t}_j$  such that  $b_j$  does not occur in a cycle with p in the dependency graph  $(1 \le j \le n)$ .

This ensures that arguments cannot grow arbitrarily through the application of recursive rules.

# Playability / Winnability

A game is <u>playable</u> if and only if every player has at least one legal move in every non-terminal state.

(Note that in chess, if a player cannot move, it is a stalemate. Fortunately, this is a terminal state.)

In GGP, every game should be playable.

A game is <u>strongly winnable</u> if and only if, for some player, there is a sequence of individual moves of that player that leads to a terminating goal state for that player.

A game is <u>weakly winnable</u> if and only if, for every player, there is a sequence of joint moves of the players that leads to a terminating goal state for that player.

In GGP, every game should be weakly winnable, and all single player games should be strongly winnable.

## Knowledge Interchange Format

Knowledge Interchange Format is a standard for programmatic exchange of knowledge represented in relational logic.

Syntax is prefix version of standard syntax. Some operators are renamed: not, and, or. Case-independent. Variables are prefixed with ?.

 $r(X,Y) <= p(X,Y) \land \neg q(Y)$ 

(<= (r ?x ?y) (and (p ?x ?y) (not (q ?y))))
or, equivalently,</pre>

(<= (r ?x ?y) (p ?x ?y) (not (q ?y)))

Semantics is the same.

#### Tic Tac Toe in KIF Notation

```
(role xplayer)
(role(oplayer)
(init (cell 1 1 b))
(init (cell 1 2 b))
(init (cell 1 3 b))
(init (cell 2 1 b))
(init (cell 2 2 b))
(init (cell 2 3 b))
(init (cell 3 1 b))
(init (cell 3 2 b))
(init (cell 3 3 b))
(init (control xplayer))
(<= (next (cell ?m ?n x))</pre>
    (does xplayer (mark ?m ?n))
(<= (next (cell ?m ?n o))</pre>
    (does oplayer (mark ?m ?n))
(<= (next (cell ?m ?n ?w))</pre>
    (true (cell ?m ?n ?w))
    (does ?p (mark ?j ?k))
    (or (distinct ?m ?j)
         (distinct ?n ?k)))
(<= (next (control xplayer))</pre>
    (true (control oplayer)))
(<= (next (control oplayer))</pre>
    (true (control xplayer)))
```

```
(<= (legal ?p (mark ?m ?n))</pre>
    (true (cell ?m ?n b))
    (true (control ?p)))
(<= (legal xplayer noop)</pre>
    (true (control oplayer)))
(<= (legal oplayer noop)</pre>
    (true (control xplayer)))
(<= (row ?m ?w)
     (true (cell ?m 1 ?w))
     (true (cell ?m 2 ?w))
     (true (cell ?m 3 ?w)))
(<= (column ?n ?w)</pre>
     (true (cell 1 ?n ?w))
     (true (cell 2 ?n ?w))
     (true (cell 3 ?n ?w)))
(<= (diagonal ?w)</pre>
     (true (cell 1 1 ?w))
     (true (cell 2 2 ?w))
     (true (cell 3 3 ?w)))
(<= (diagonal ?x)</pre>
     (true (cell 1 3 ?w))
     (true (cell 2 2 ?w))
     (true (cell 3 1 ?w)))
```

```
(<= (line ?w) (row ?m ?w))</pre>
(<= (line ?w) (column ?n ?w))</pre>
(<= (line ?w) (diagonal ?w))</pre>
(<= open
     (true (cell ?m ?n b)))
(<= terminal</pre>
     (or (line x) (line o)))
(<= terminal
     (not open))
(<= (goal xplayer 100)</pre>
     (line x))
(<= (qoal xplayer 50)</pre>
     draw)
(<= (qoal xplayer 0)</pre>
     (line o))
(<= (goal oplayer 100)</pre>
     (line o))
(<= (goal oplayer 50)</pre>
     draw)
(<= (goal oplayer 0)</pre>
     (line x))
```

# **Playing Games**

# www.general-game-playing.de/downloads.html

Downloads - General Game Playing

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# **Communication Protocol**

- - Role: the name of the role you are playing (e.g. xplayer or oplayer)
  - Game description: the axioms describing the game
  - Start/play clock: how much time you have before the game begins/per turn
- Manager sends PLAY message to players
   (PLAY <MATCH ID> <PRIOR MOVES>)

Prior moves is a list of moves, one per player

- The order is the same as the order of roles in the game description
- e.g. ((mark 1 1) noop)
- Special case: for the first turn, prior moves is nil
- Players send back a message of the form MOVE, e.g. (mark 3 2)
- When the previous turn ended the game, Manager sends a STOP message (STOP <MATCH ID> <PRIOR MOVES>)

### GameControllerApp



# Implementing a Basic General Game Player:

# **Blind Search**

# Single-Player Games: A Simple Example



Pressing button *a* toggles *p*.

Pressing button b interchanges p and q.

Initially, *p* and *q* are off. Goal: *p* and *q* are on.

#### **Game Description**

(role robot)

#### Legality

(legal robot a)
(legal robot b)

#### Update

(<= (next(p) (does robot a) (not (true p))) (<= (next(q) (does robot a) (true q)) (<= (next(p) (does robot b) (true q)) (<= (next(q) (does robot b) (true p))</pre>

Termination and Goal

```
(<= terminal (true p) (true q))
(<= (goal robot 100) (true p) (true q))</pre>
```

# Solving Single-Player Games = Planning

- Initial state
  - {} (since there is no rule for **init** in this game)
- Actions
  - a Preconditions: none
     Effects: toggles truth-value of p
  - b Preconditions: none
    - Effects: interchanges truth-values of  $\ {\rm p}\$  and  $\ {\rm q}\$
- Goal
  - p ^ q

#### State Transition System



# Single-Player, Perfect-Information Games

Many single-player games can be solved using standard search techniques

- Iterative deepening
- Bidirectional search

Special techniques

• Constraint solving (suitable for Sudoku, Gene Sequencing and the like)

### Multi-Player Games: Game Tree Search



### How to Deal With Simultaneous Moves





#### Minimax With $\alpha$ - $\beta$ -Heuristics



#### Stochastic Search (1)



Minimax Search



Monte Carlo Tree Search (random simulations)

### Stochastic Search (2)

Value of move = Average score returned by simulation



# Stochastic Search (3): Confidence Bounds

- Play one random game for each move
- For next simulation choose move

